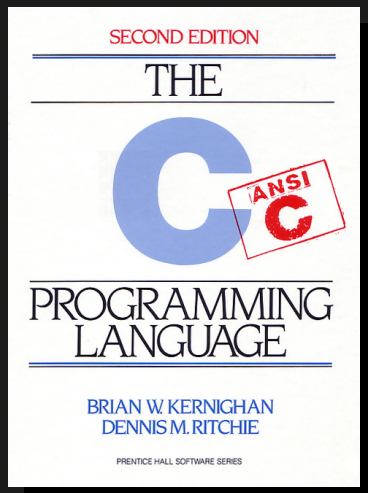
More Number Representation, C Programming Language Overview



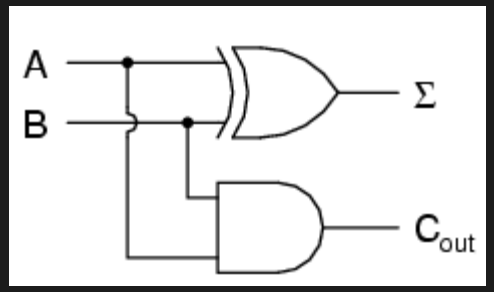
Math 230

Assembly Language Programming (Computer Organization)

Lecture 3

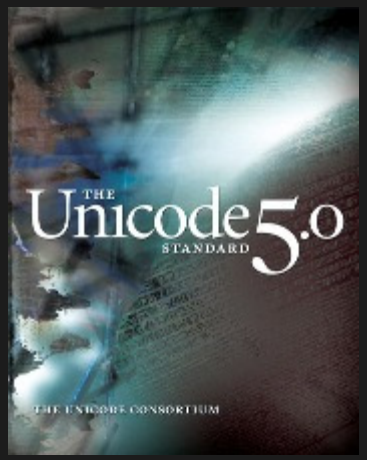
Operations with Numbers

- What things can we do with numbers?
  - Add
  - Subtract
  - Multiply
  - Divide
  - Compare
- Addition
  - can easily build a circuit to do it!
- Subtraction
  - can we use the circuitry for addition?
- Comparison
  - How can you tell if one number is larger than another?



Bits Can Represent Anything

- Colors, music, text,video, texture, smell
  - We can quantize anything
- Characters
  - 26 letters => 5 bits
  - upper + lowercase+ punctuations => 7 bits (ASCII)
  - to cover all world's various alphabets
    - => 8, 16, 32 bits ("Unicode")
- Logical Values
  - 0 = false
  - 1 = true
- Colors
  - using red, blue, and yellow we can represent a wide range of colors using paint
- Locations/Addresses/Commands
  - Memorize: N bits ← → 2^N things



Negative Numbers

- So far, only positive numbers?
- Consider:
  - Borrow most significant bit and call it a sign bit

How to Represent Negative Numbers?

- So far, unsigned numbers
- Obvious solution: define leftmost bit to be sign!
  - 0 ⇒ +, 1 ⇒ –
  - Rest of bits can be numerical value of number
- Representation called sign and magnitude
- MIPS uses 32-bit integers. $+1_{ten}$ would be:

0000 0000 0000 0000 0000 0000 0000 0001

- And -1_{ten} in sign and magnitude would be:

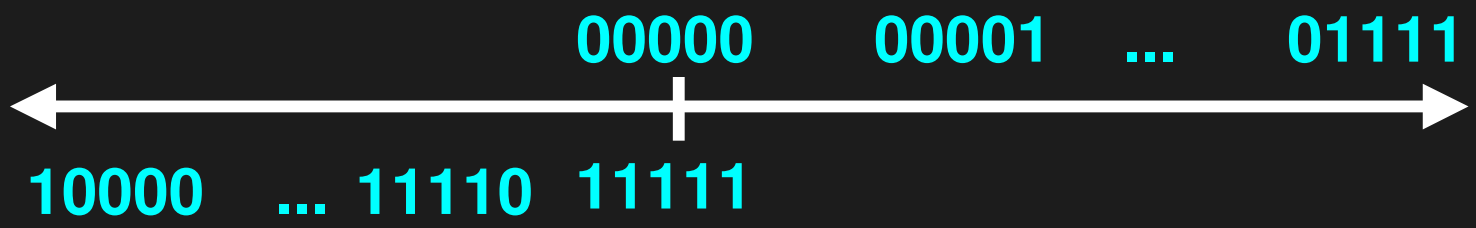
1000 0000 0000 0000 0000 0000 0000 0001

Shortcomings of sign and magnitude?

- Arithmetic circuit complicated
  - Special steps depending whether signs are the same or not
- Also, two zeros
  - $0x00000000 = +0_{ten}$
  - $0x80000000 = -0_{ten}$
  - What would two 0s mean for programming?
- Therefore sign and magnitude abandoned

Another try: complement the bits

- Example: $7_{10} = 00111_2$ $-7_{10} = 11000_2$
- Called One's Complement
- Note: positive numbers have leading 0s, negative numbers have leadings 1s.



- What is -00000 ? Answer: 11111
- How many positive numbers in N bits?
- How many negative numbers?

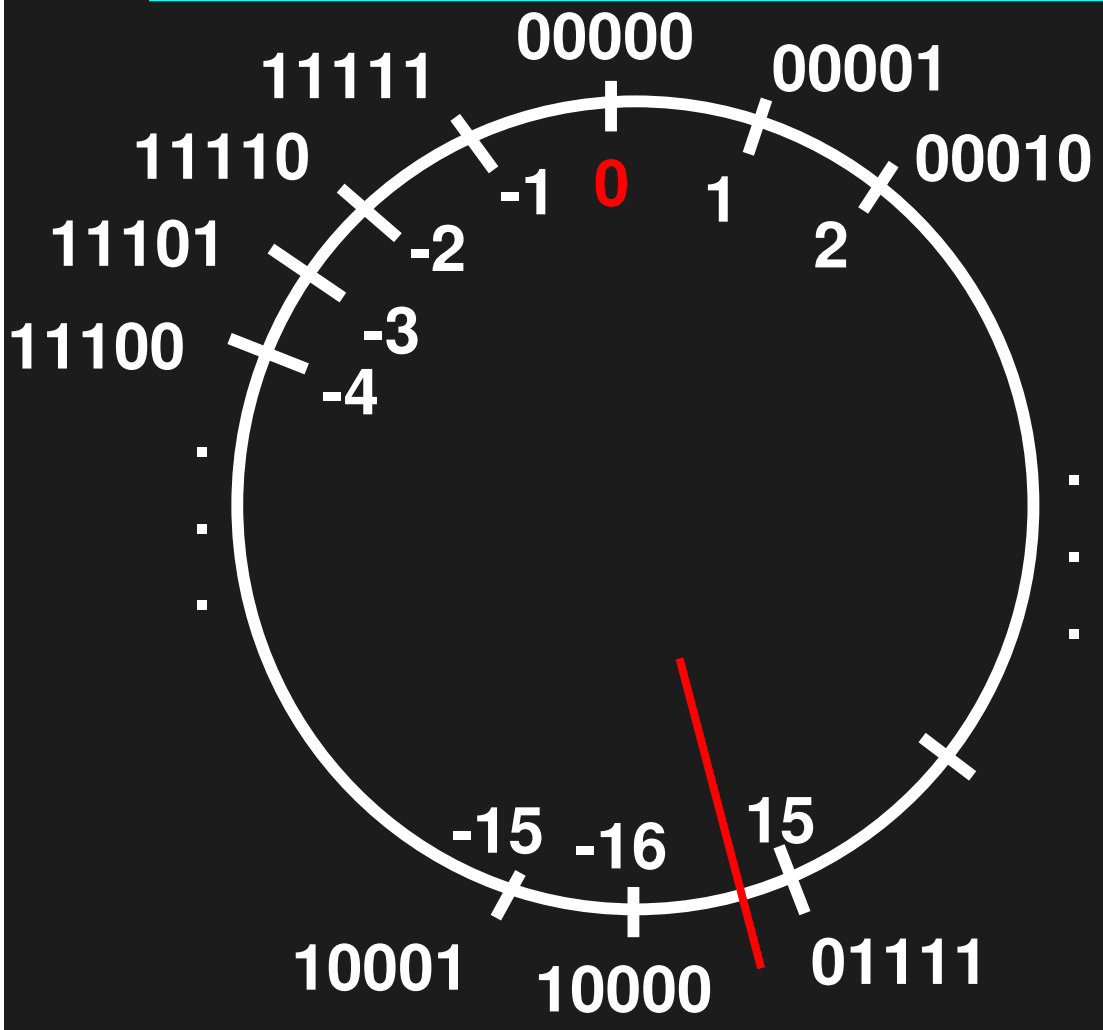
Shortcomings of One's complement?

- Arithmetic still somewhat complicated.
- Still two zeros
  - `0x00000000` = $+0_{ten}$
  - `0xFFFFFFFF` = -0_{ten}
- Although used for awhile on some computer products, one's complement was eventually abandoned because another solution was better.

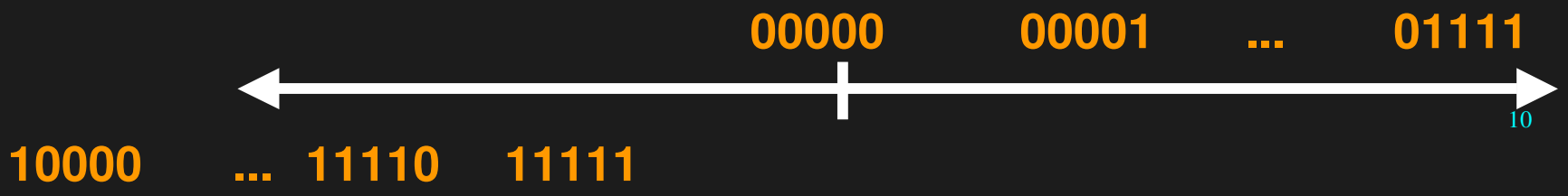
Standard Negative Number Representation

- What is result for unsigned numbers if tried to subtract large number from a small one?
  - Would try to borrow from string of leading 0s, so result would have a string of leading 1s
    - 3 - 4 $\Rightarrow$ 00…0011 – 00…0100 = 11…1111
  - With no obvious better alternative, pick representation that made the hardware simple
  - As with sign and magnitude, leading 0s $\Rightarrow$ positive, leading 1s $\Rightarrow$ negative
    - 000000...xxx is ≥ 0, 111111...xxx is < 0
    - except 1…1111 is -1, not -0 (as in sign & mag.)
- This representation is Two's Complement

2's Complement Number "line": N = 5



- 2^{N-1} non-negatives
- 2^{N-1} negatives
- one zero
- how many positives?



Two's Complement for N=32

$0000 \ldots 0000\ 0000\ 0000\ 0000_{two}$ =	0_{ten}
$0000 \ldots 0000\ 0000\ 0000\ 0001_{two}$ =	1_{ten}
$0000 \ldots 0000\ 0000\ 0000\ 0010_{two}$ =	2_{ten}
. . .	
$0111 \ldots 1111\ 1111\ 1111\ 1101_{two}$ =	$2{,}147{,}483{,}645_{ten}$
$0111 \ldots 1111\ 1111\ 1111\ 1110_{two}$ =	$2{,}147{,}483{,}646_{ten}$
$0111 \ldots 1111\ 1111\ 1111\ 1111_{two}$ =	$2{,}147{,}483{,}647_{ten}$
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. . .	
$1111 \ldots 1111\ 1111\ 1111\ 1101_{two}$ =	-3_{ten}
$1111 \ldots 1111\ 1111\ 1111\ 1110_{two}$ =	-2_{ten}
$1111 \ldots 1111\ 1111\ 1111\ 1111_{two}$ =	-1_{ten}

- One zero; 1st bit called sign bit
- 1 "extra" negative:no positive $2{,}147{,}483{,}648_{ten}$

Two's Complement Formula

- Can represent positive and negative numbers in terms of the bit value times a power of 2:

  $d_{31} \times -(2^{31}) + d_{30} \times 2^{30} + \ldots + d_2 \times 2^2 + d_1 \times 2^1 + d_0 \times 2^0$

- Example: 1101_{two}

  $= 1 \times -(2^3) + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$

  $= -2^3 + 2^2 + 0 + 2^0$

  $= -8 + 4 + 0 + 1$

  $= -8 + 5$

  $= -3_{ten}$

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Two's Complement shortcut: Negation

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- Change every 0 to 1 and 1 to 0 (invert or complement), then add 1 to the result

Two's comp. shortcut: Sign extension

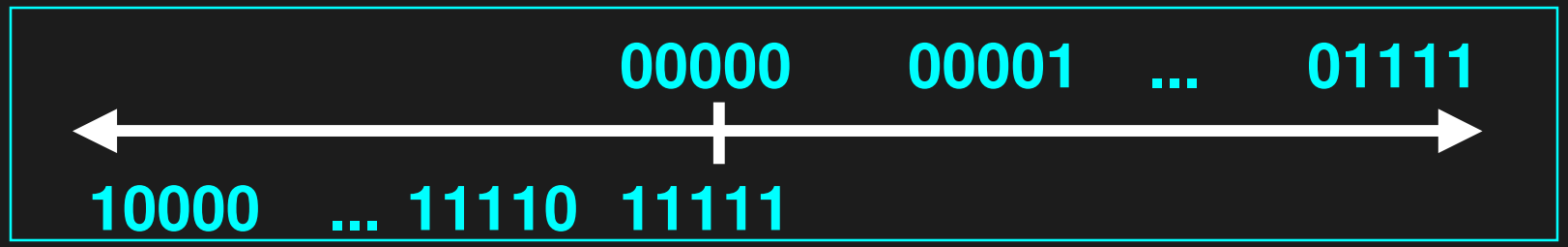
- Convert 2's complement number rep. using n bits to more than n bits
- Simply replicate the most significant bit (sign bit) of smaller to fill new bits
  - 2's comp. positive number has infinite 0s
  - 2's comp. negative number has infinite 1s
  - Binary representation hides leading bits; sign extension restores some of them
  - 16-bit -4_{ten} to 32-bit:

$$1111\ 1111\ 1111\ 1100_{two}$$

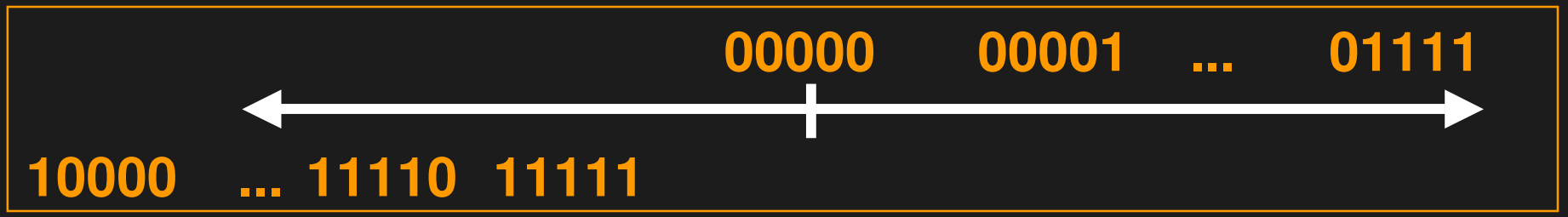
$$1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1100_{two}$$

Number summary…

- We represent "things" in computers as particular bit patterns: N bits ⇒ 2^N
- Decimal for human calculations, binary for computers, hex to write binary more easily
- 1's complement - mostly abandoned



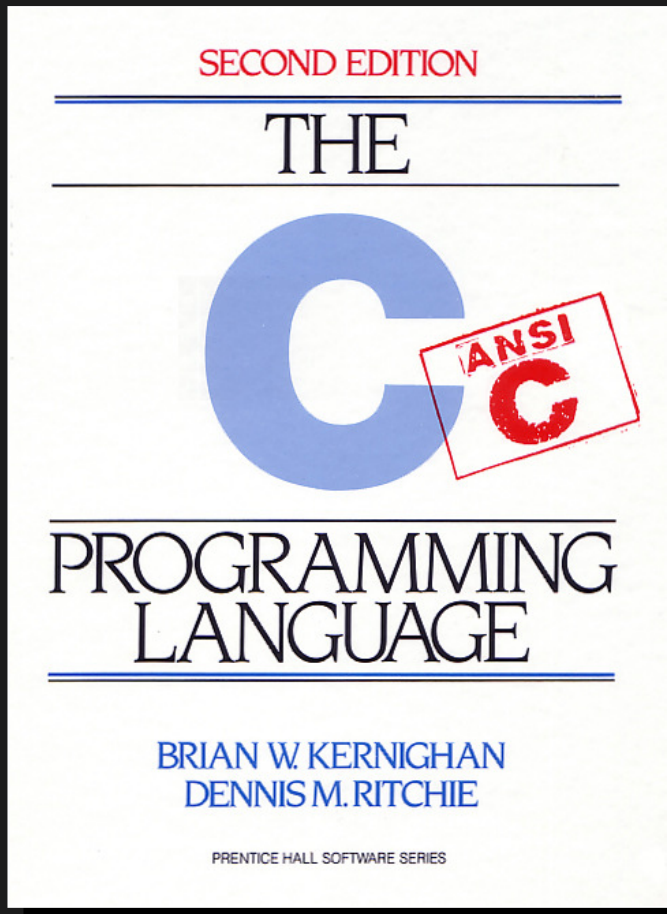
- 2's complement universal in computing: cannot avoid, so learn



- Overflow: numbers ∞; computers finite,errors!

C Overview

- Systems software
  - Unix, Linux, MacOSX
- Small Language
- Standardized C Preprocessor for
  - macro def
  - source code file inclusion
  - conditional compilation
- gcc can cross compile to MIPS et al



ANSI C

- K&R was the informal “spec” for many years
  - K&R C
  - 2nd edition covers ANSI C std
- ANSI: Superset of K&R *C*
- After K&R published
  - **`void`** functions
  - functions returning **`struct`** or **`union`**
  - assignment for **`struct`**
  - **`const`** qualifier

```
int main(argc, argv)
 int argc;
 char* argv[];
 {  .  .  .   }
```

```
int main(int argc, char* argv[])
{  .  .  .}
```

C99

- inline functions
  - eliminate function call; expand inline
- variables declared anywhere
- New data types
  - long long int
  - boolean data type
  - complex types for complex numbers
- Header files added, notably
  - `<stdbool.h>`
  - `<tgmath.h>` type generic math functions
- gcc not fully compliant (http://gcc.gnu.org/c99status.html)

Standard Headers

<assert.h>
<complex.h>
<ctype.h>
<errno.h>
<fenv.h>
<float.h>
<inttypes.h>
<iso646.h>
<limits.h>
<locale.h>
<math.h>
<setjmp.h>
<signal.h>
<stdarg.h>
<stdbool.h>
<stddef.h>
<stdint.h>
<stdio.h>
<stdlib.h>
<string.h>
<tgmath.h>
<time.h>
<wchar.h>
<wctype.h>

C vs Java

C	*Java*
Relatively Fast	*Relatively Slow*
Procedural	*OOP*
Platform Dependent	*Platform Independent*
Arrays initialize to garbage	*Arrays initialize to zero*
Small libraries	*Huge libraries*
Small executable	*Larger runnable files*
x.c => x.exe or x.out	*x.java => x.class*
Preprocessor	*No preprocessor*
No memory management	*Garbage Collections*
Pointers	*No global variables*
	Variable Declarations

C Syntax: variable declaration

- Similar to Java
- ANSI requires declarations to go at start of block
  - Java allows anywhere

C Syntax

- Java has booleans to represent true/false
- C has nonzeroes which evaluate to "true"
  - false
    - if(0)… //always false
    - if (NULL) …//always false
    - if('0') …//always true

C Syntax: Flow Control

- Decision structures of C almost identical to Java.
  - if-else
  - switch
  - while and for
  - do- while

Memory: Conceptual

- Think of memory as one HUGE array
  - we will talk about how a "2D" array maps onto one linear space
- Each cell can be "*addresed*"
  - Address signed or unsigned?
- Commonly work with reference-type variables, i.e., pointers

Pointers

Memory Address of a Variable

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```
char ch = 'A';
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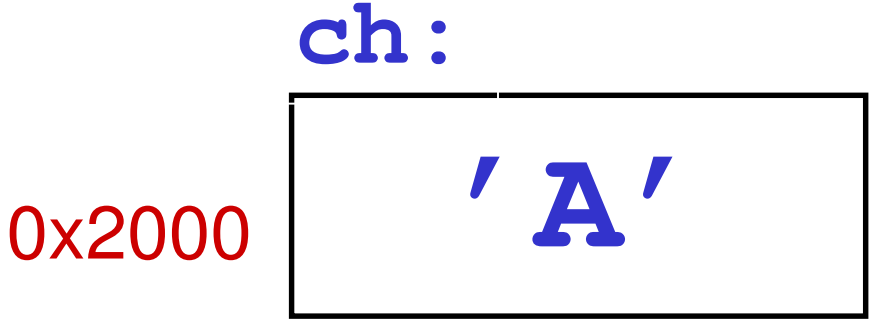

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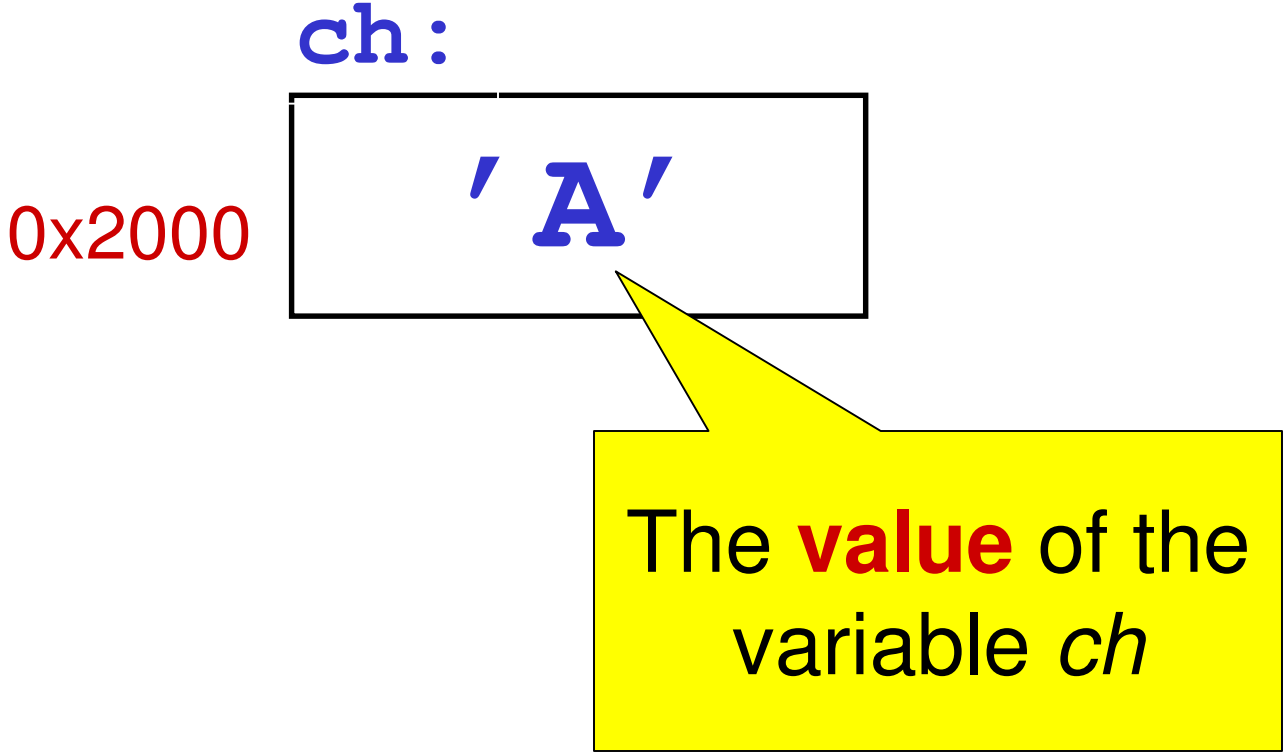

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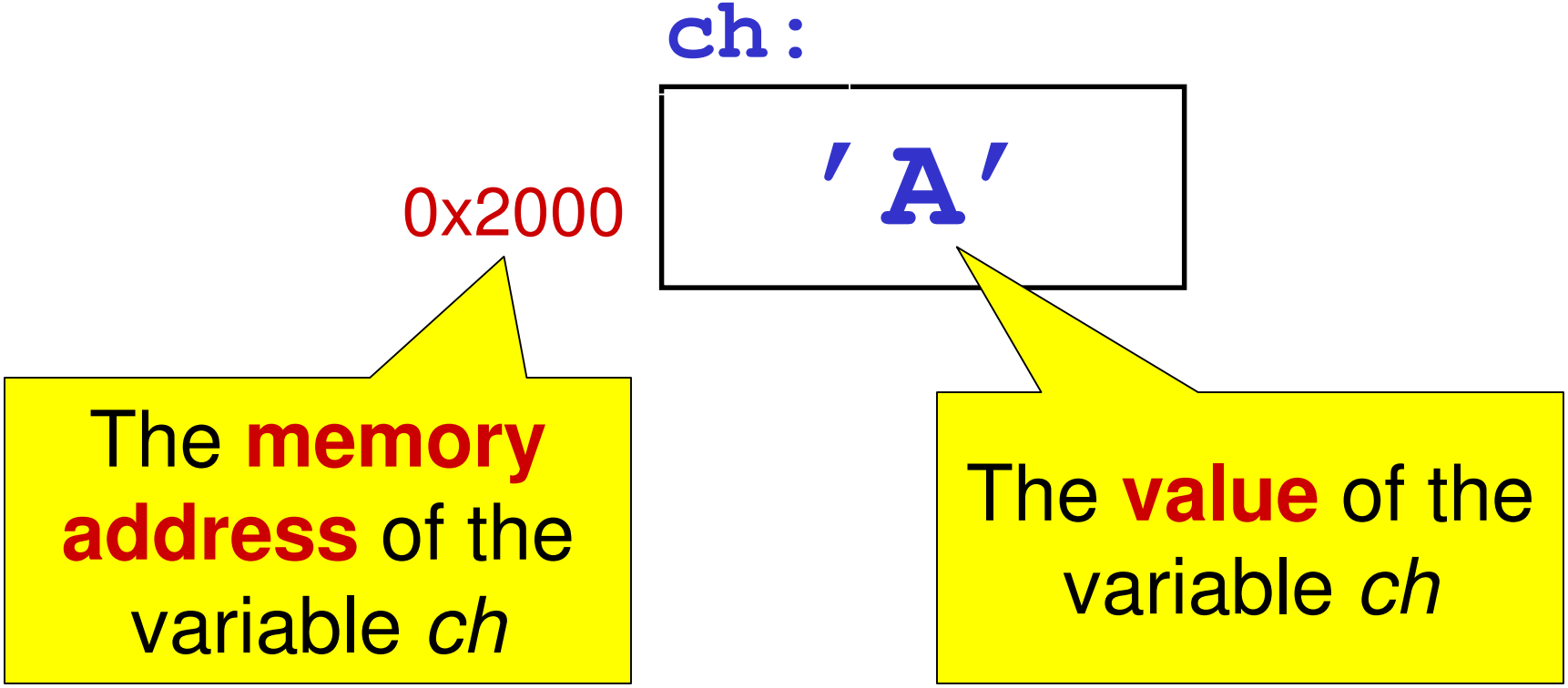

The & Operator

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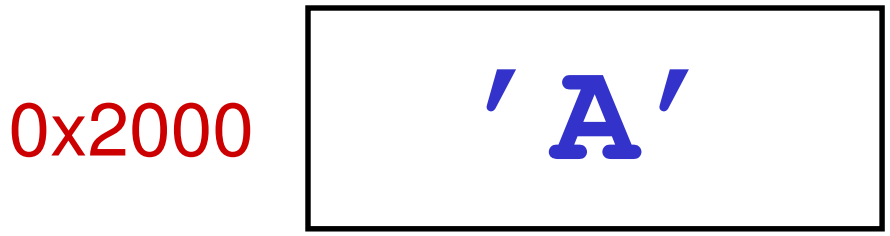

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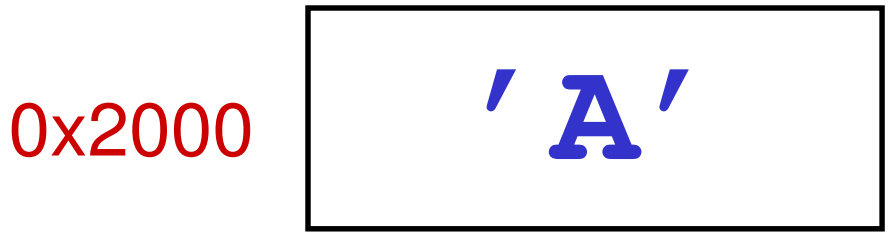

`&ch` yields the value 0x2000

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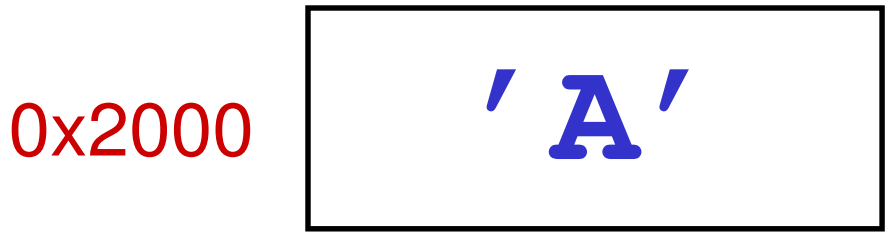

`&ch` yields the value 0x2000

- Also known as the "address operator"

Example:

```
char ch;
printf("%p", &ch);
```

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```
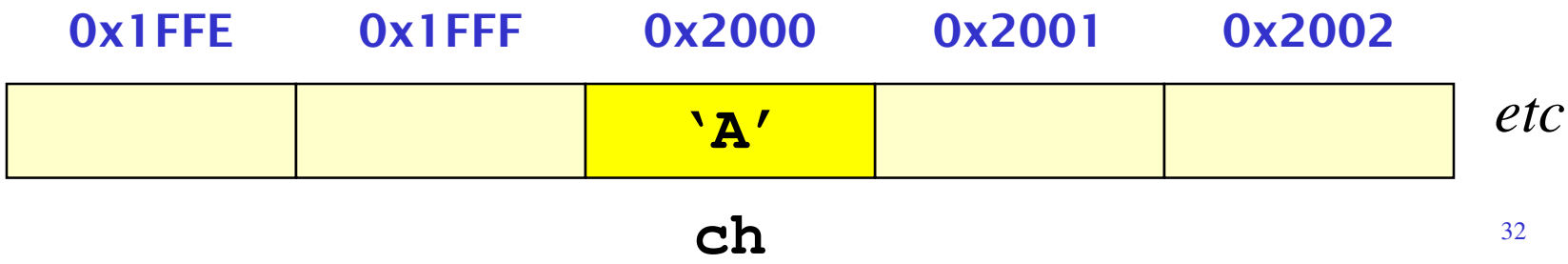
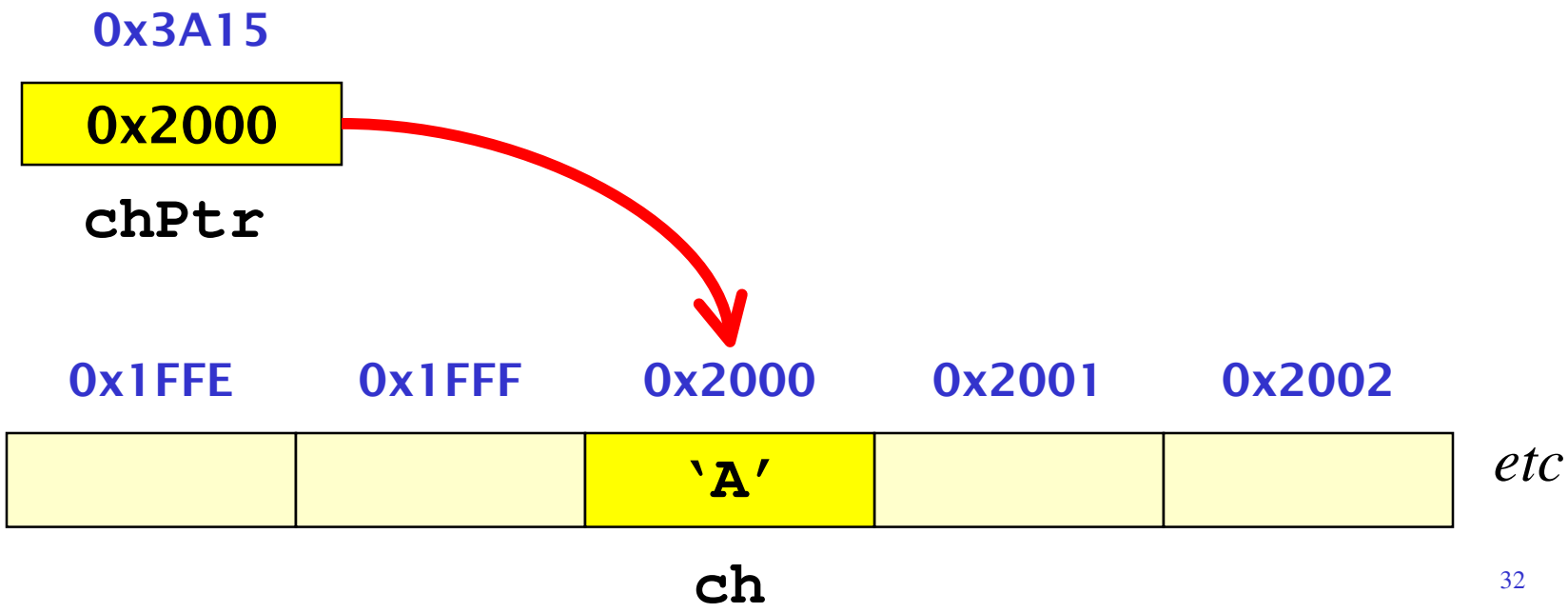

"conversion specifier" for printing a memory address

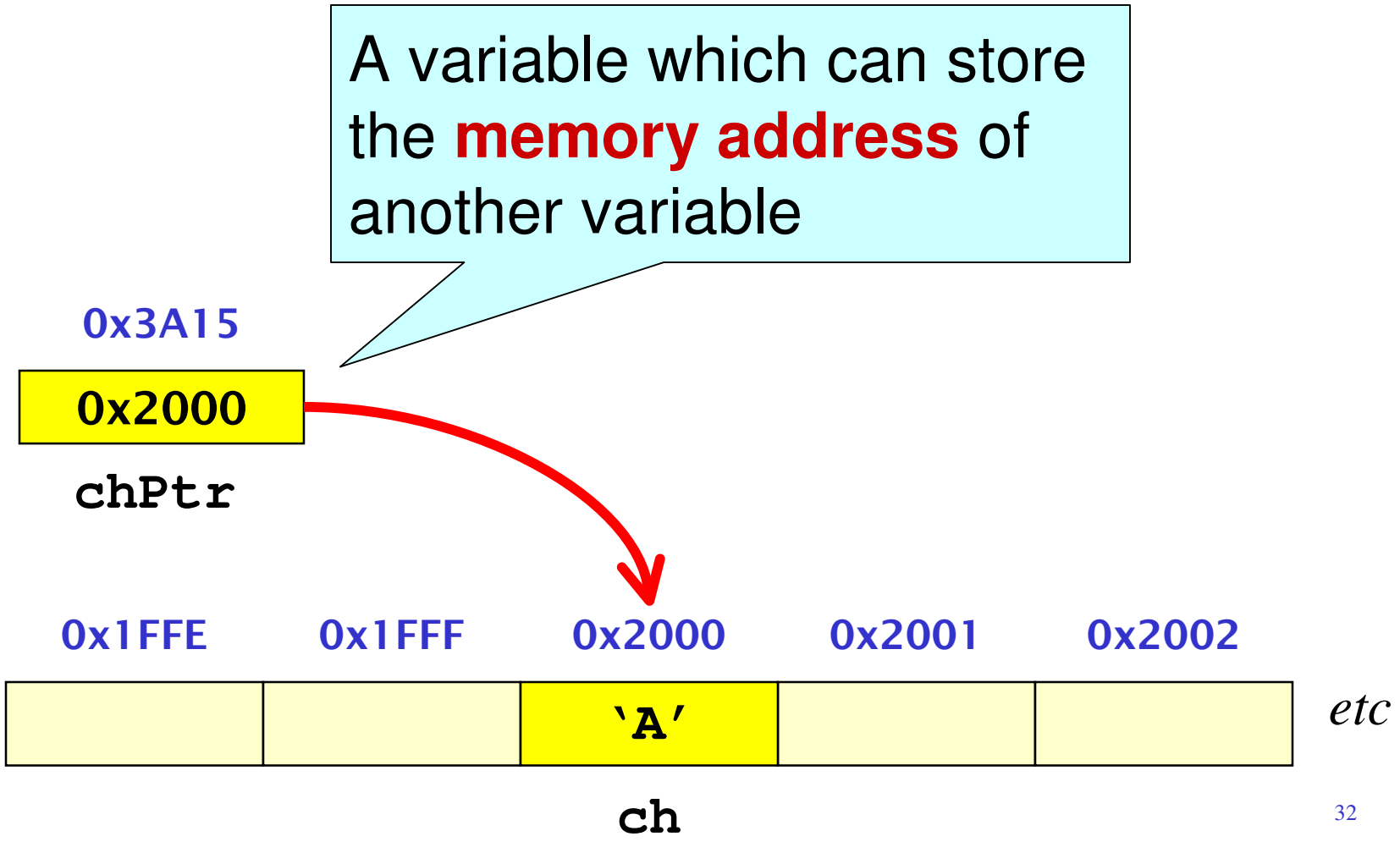
Pointers



Pointers


0x3A15
0x2000
chPtr
0x1FFE
0x1FFF
0x2000
0x2001
0x2002
'A'
etc
ch

Pointers


A variable which can store the memory address of another variable
0x3A15
0x2000
chPtr
0x1FFE
0x1FFF
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etc
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Pointers

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- A pointer is a variable which…
  - Contains a memory address
  - Points to a specific data type

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- A pointer is a variable which…
  - Contains a memory address
  - Points to a specific data type
- Pointer variables are usually named ***`varPtr`***

Example:

```
char* cPtr;
```

cPtr:

0x2004

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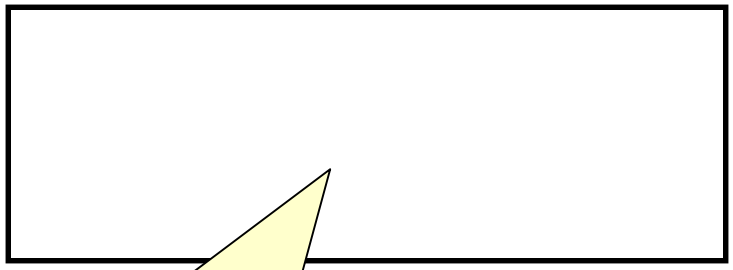
`cPtr:`

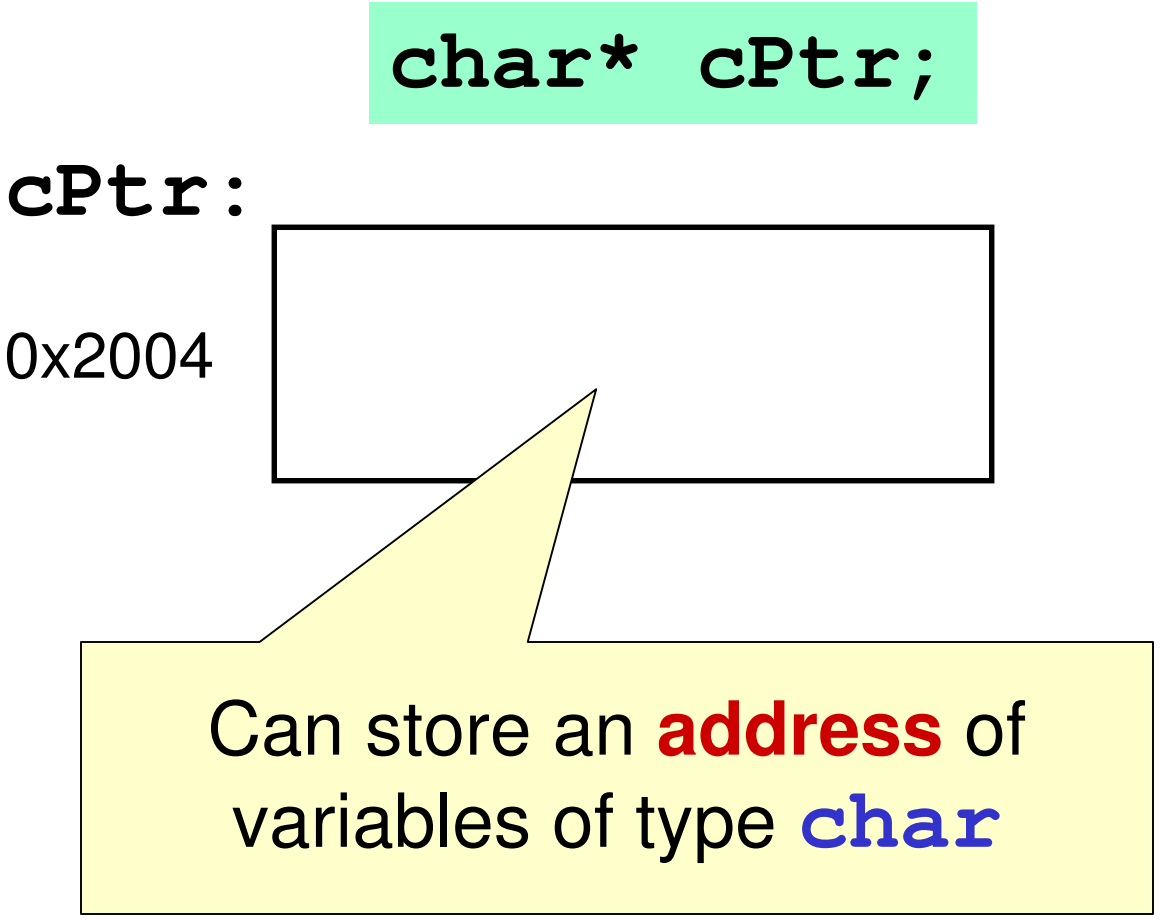
0x2004



Can store an **address** of variables of type **`char`**

Example:

```
char* cPtr;
```



- We say *cPtr* is a ***pointer*** to a character

Pointers and the & Operator

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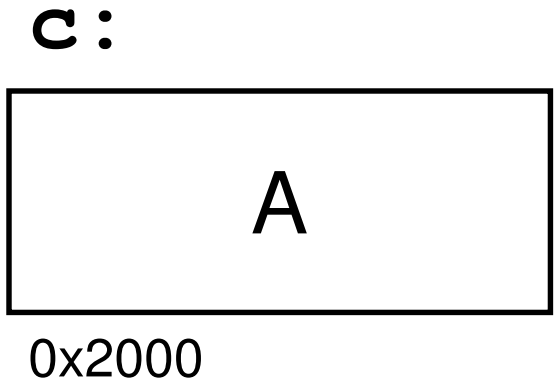

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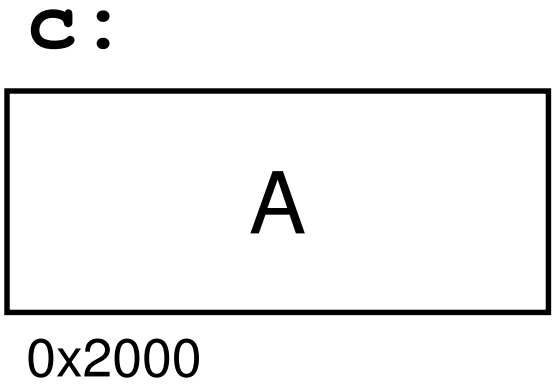

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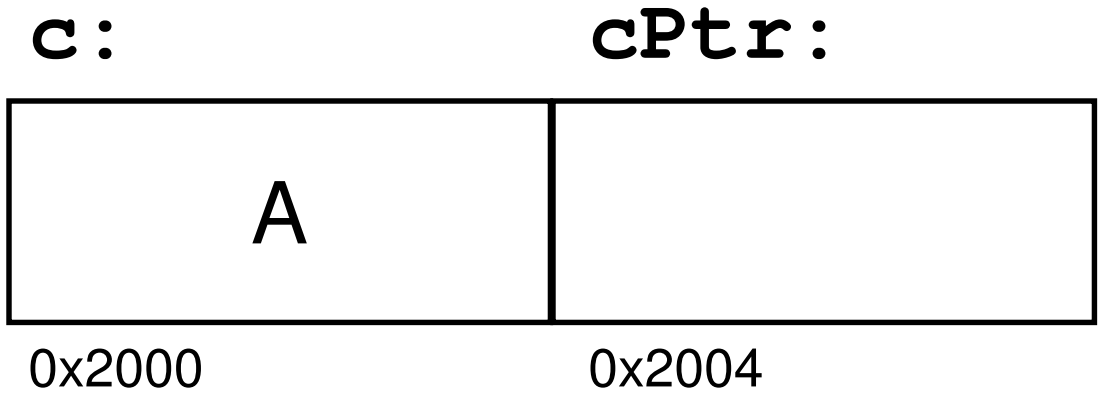

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char c = 'A';
char *cPtr;


cPtr = &c;
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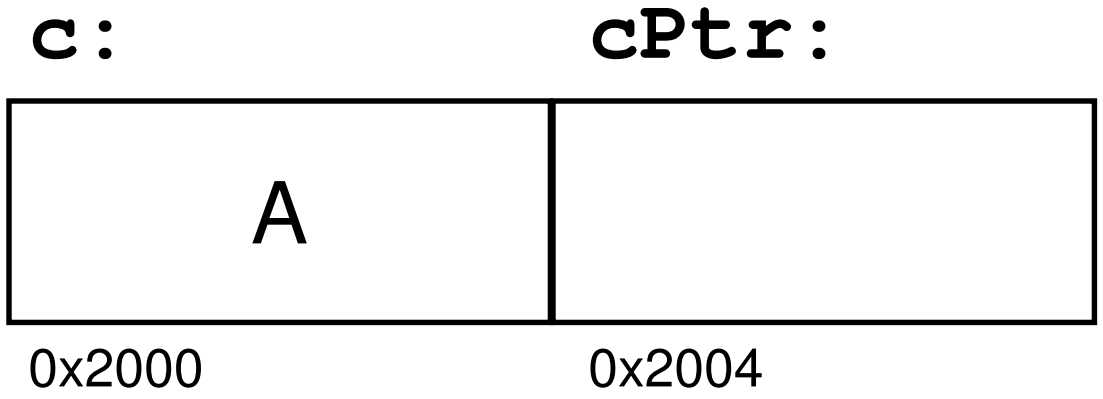

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char c = 'A';
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cPtr = &c;
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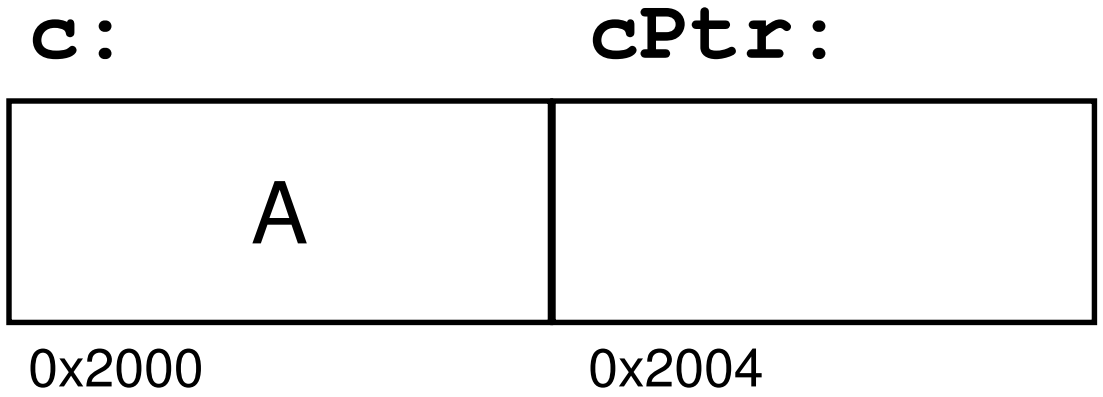
Assigns the address of c to cPtr



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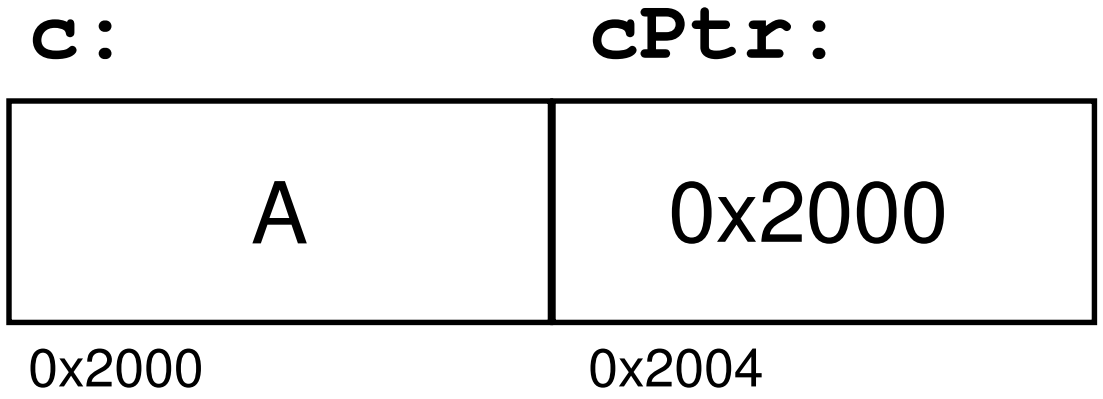

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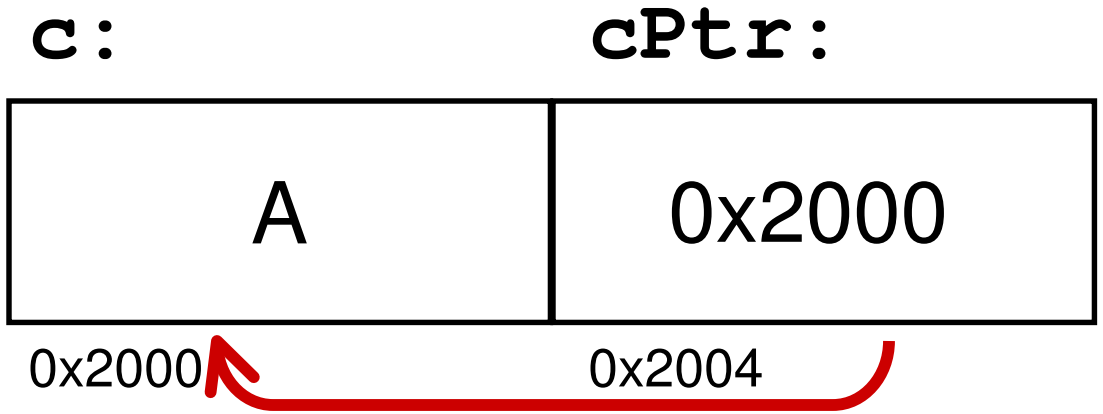

Notes on Pointers

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- We can have pointers to any data type

Example:

```
int*    numPtr;
float*   xPtr;
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- The * can be anywhere between the type and the variable

Example:

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int     *numPtr;
float * xPtr;
```

Notes on Pointers (cont)

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- You can assign the address of a variable to a "compatible" pointer using the **&** operator

Example:

```
int   aNumber;
int   *numPtr;

numPtr = &aNumber;
```

Notes on Pointers (cont)

- You can assign the address of a variable to a "compatible" pointer using the **&** operator

Example:

```
int   aNumber;
int   *numPtr;

numPtr = &aNumber;
```

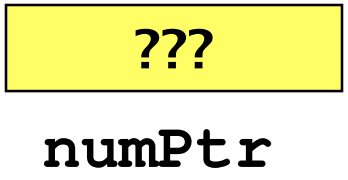
- You can print the address stored in a pointer using the **%p** conversion specifier

Example:

```
printf("%p", numPtr);
```

Notes on Pointers (cont)

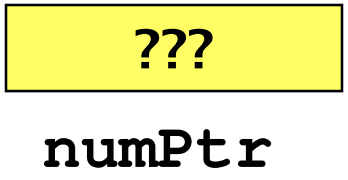
```
int   *numPtr;
```



Notes on Pointers (cont)

```
int   *numPtr;
```

Beware of pointers which are not initialized!



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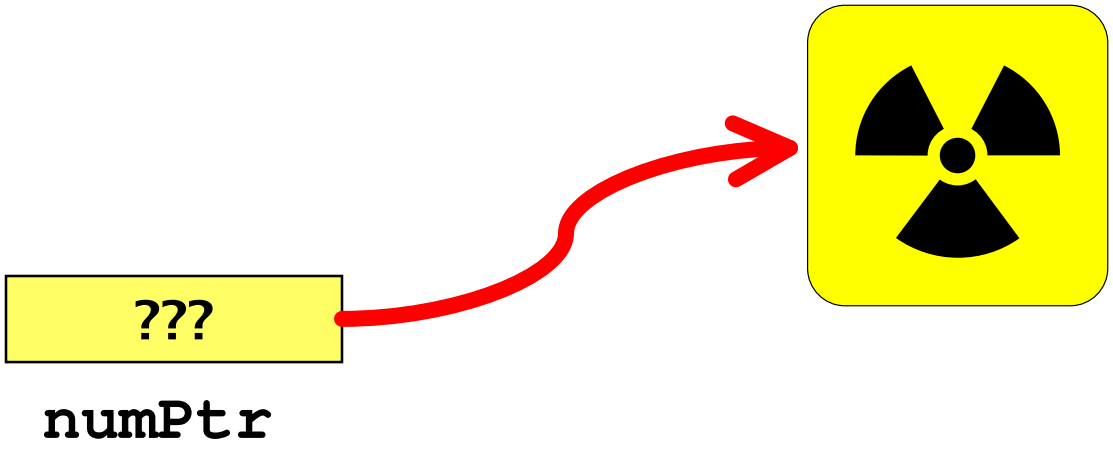
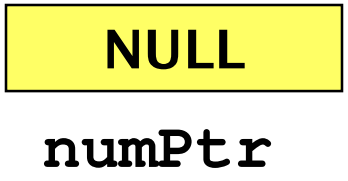


Notes on Pointers (cont)

- When declaring a pointer, it is a good idea to always initialize it to **NULL** (a special pointer constant)

```
int   *numPtr = NULL;
```



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- Should not be confused with the * in the pointer declaration

Pointers and the * Operator

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char c = 'A';
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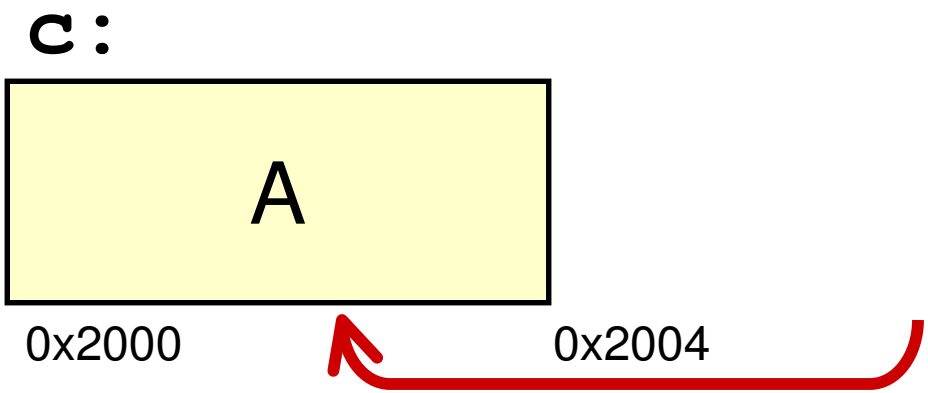

Pointers and the * Operator

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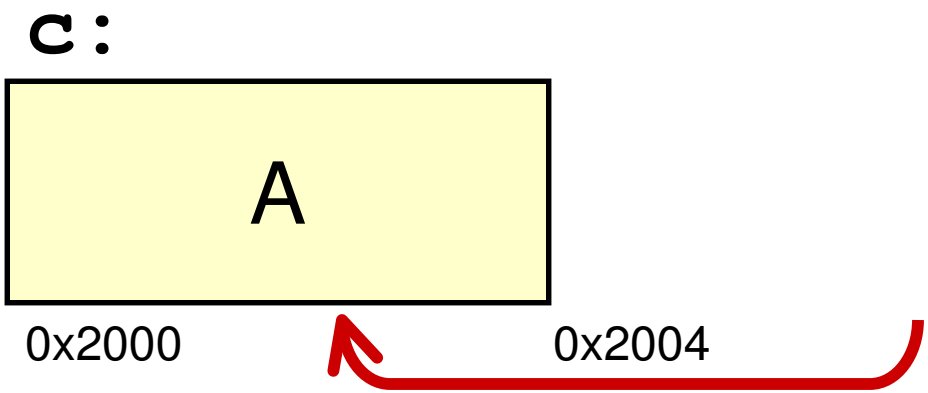
char c = 'A';
char *cPtr = NULL;
```



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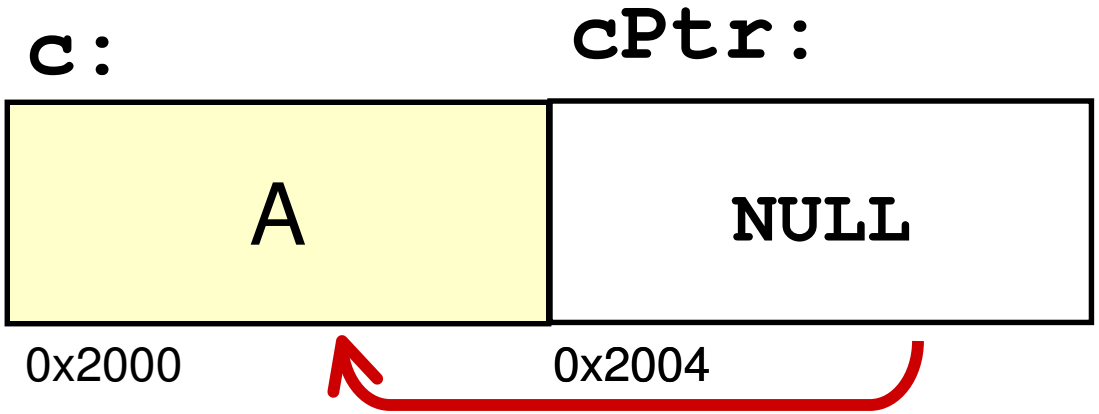

Pointers and the * Operator

Example:

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char c = 'A';
char *cPtr = NULL;


cPtr = &c;
```

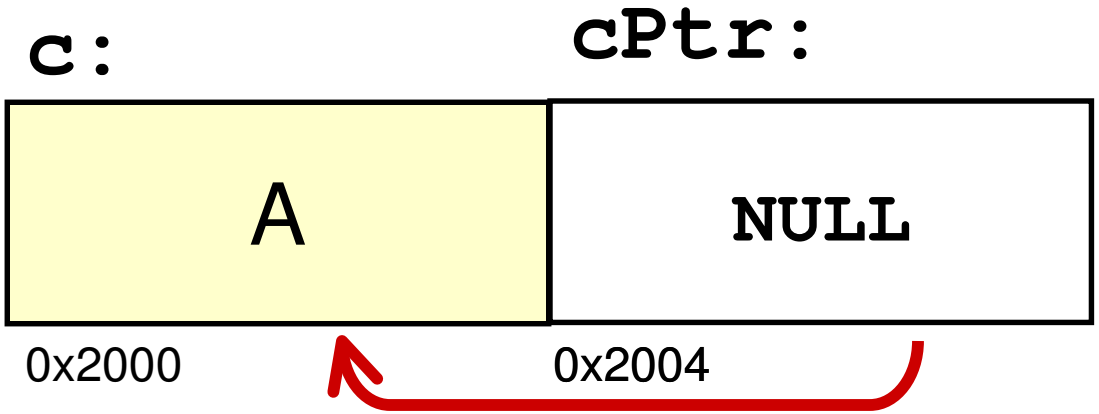

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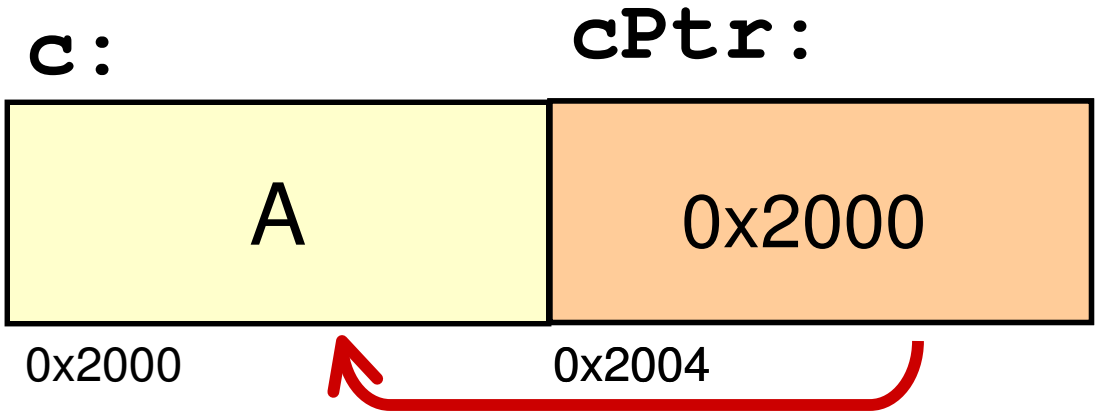
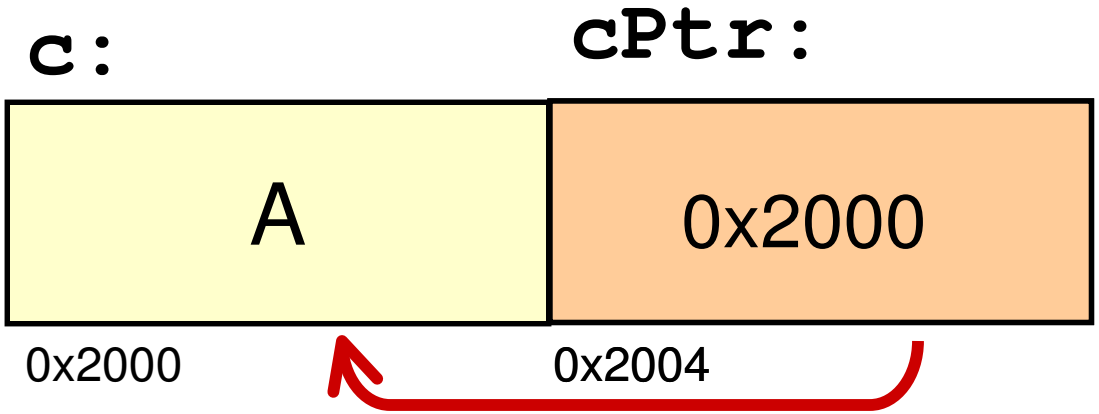


Pointers and the * Operator

Example:

```
char c = 'A';
char *cPtr = NULL;


cPtr = &c;
*cPtr = 'B';
```



Pointers and the * Operator

Example:

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char c = 'A';
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cPtr = &c;
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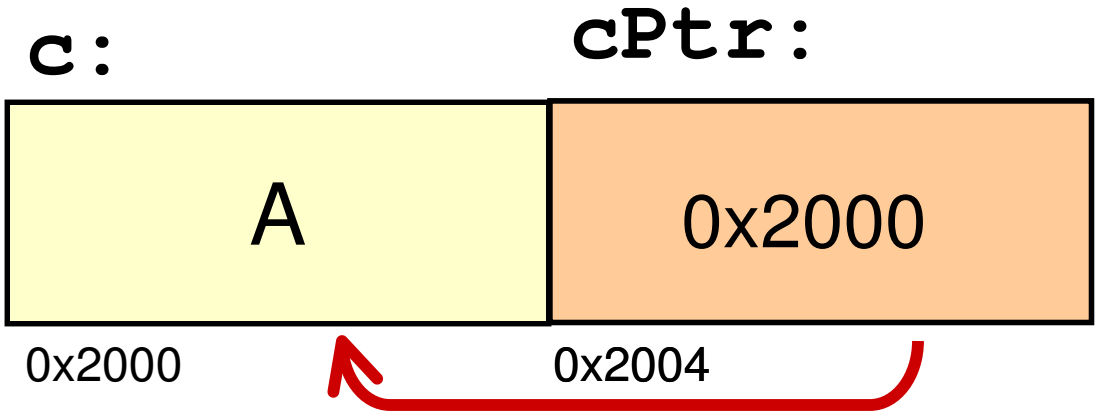
Changes the value of the variable which cPtr points to



Pointers and the * Operator

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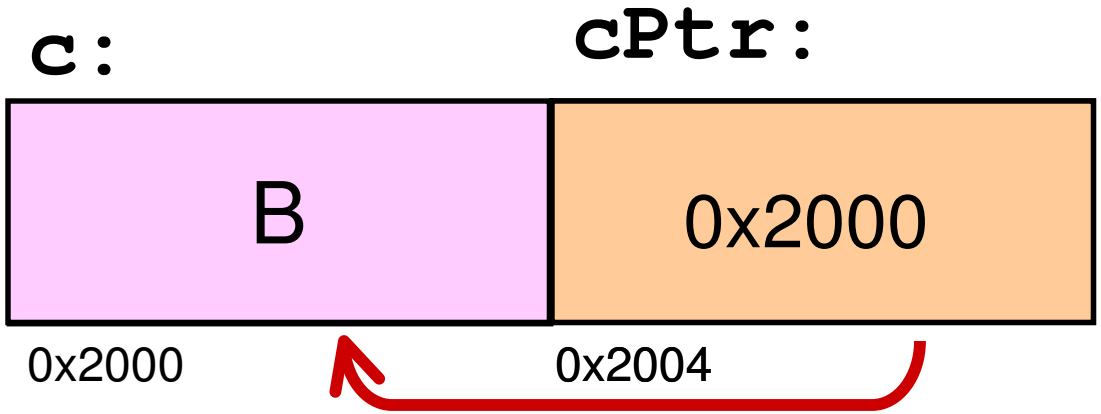

Lessons

- A pointer is simply a variable that contains an address